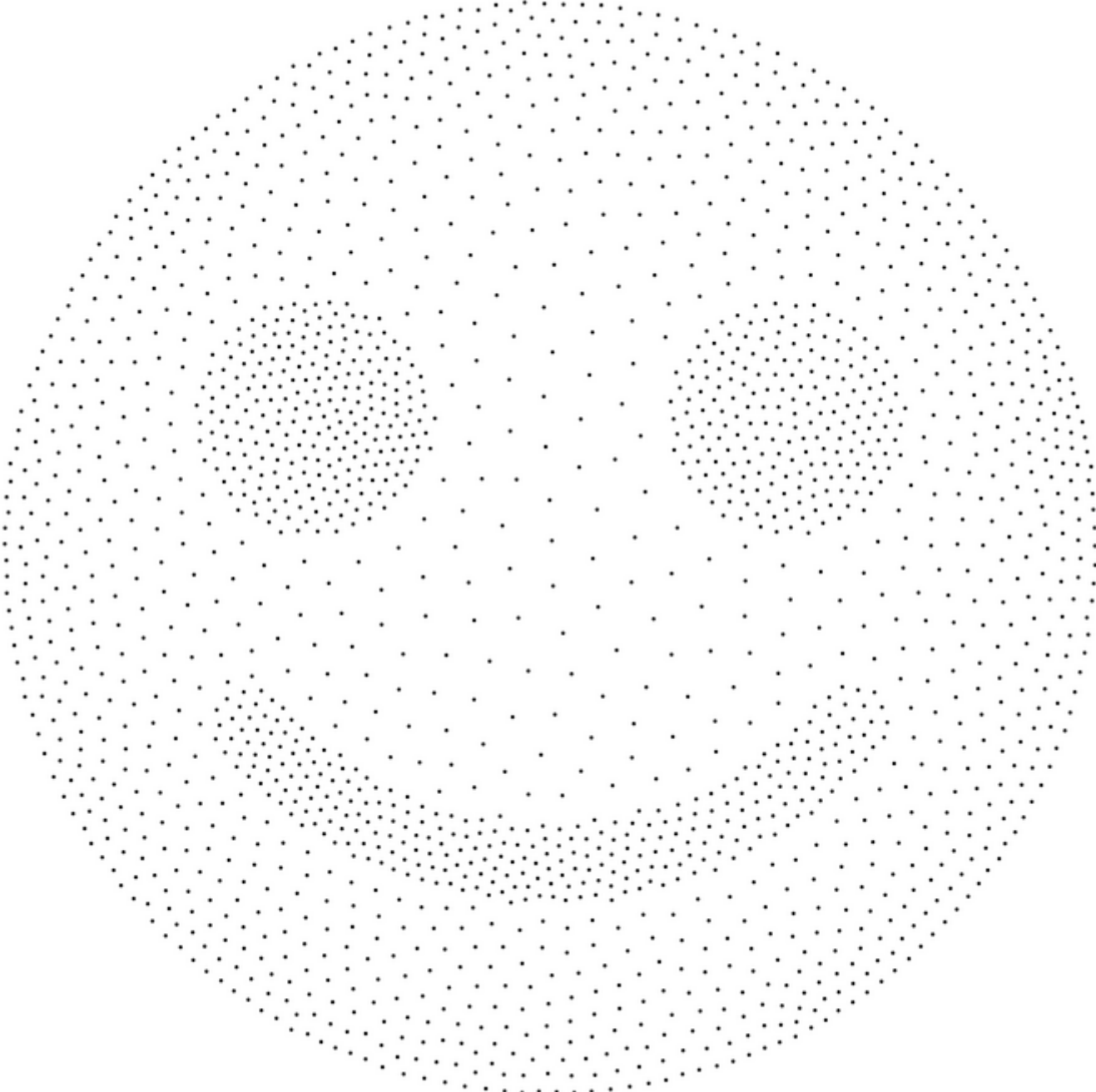


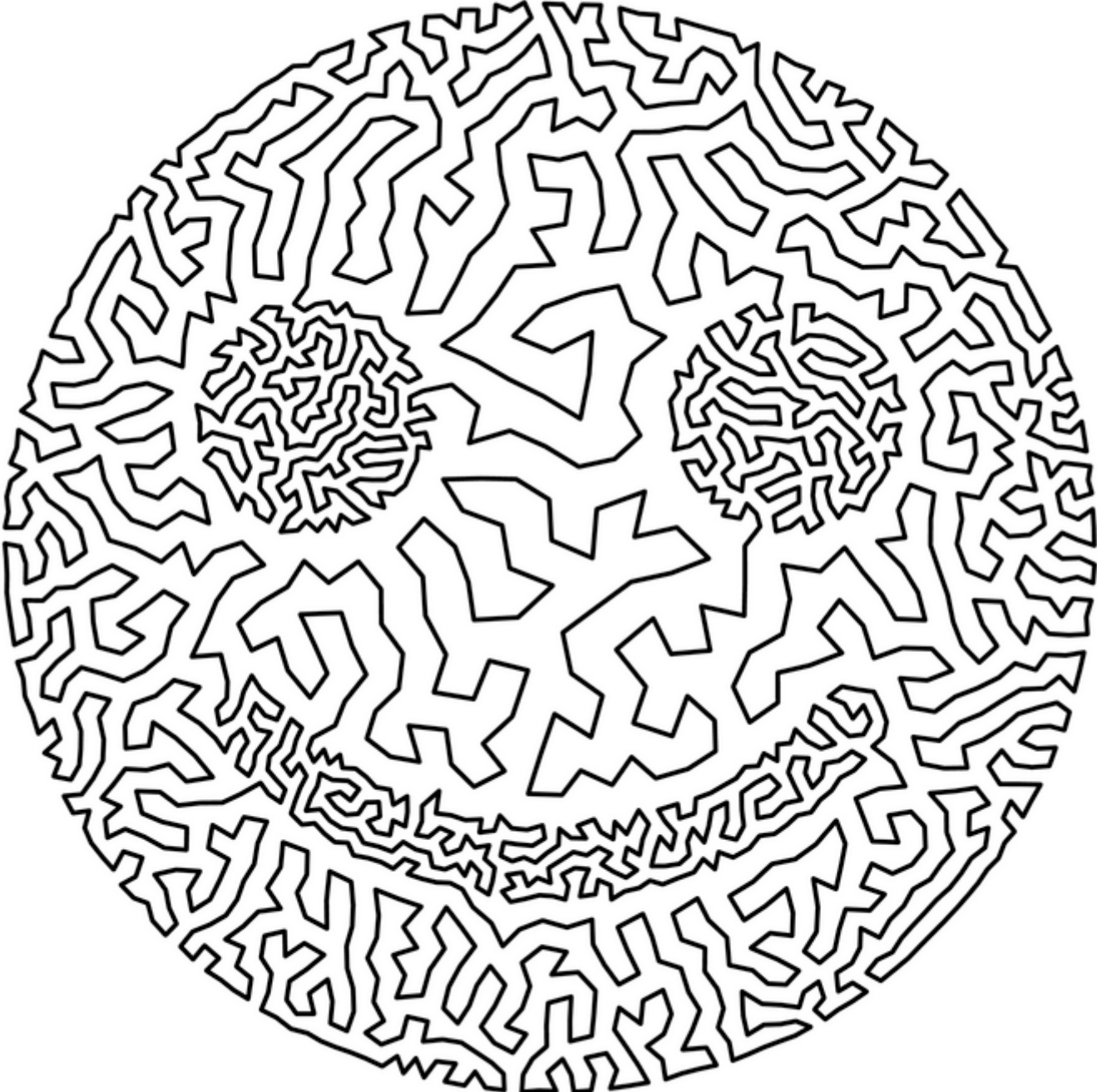
TSPortraits of Knots and Links

Robert Bosch

May 13, 2009



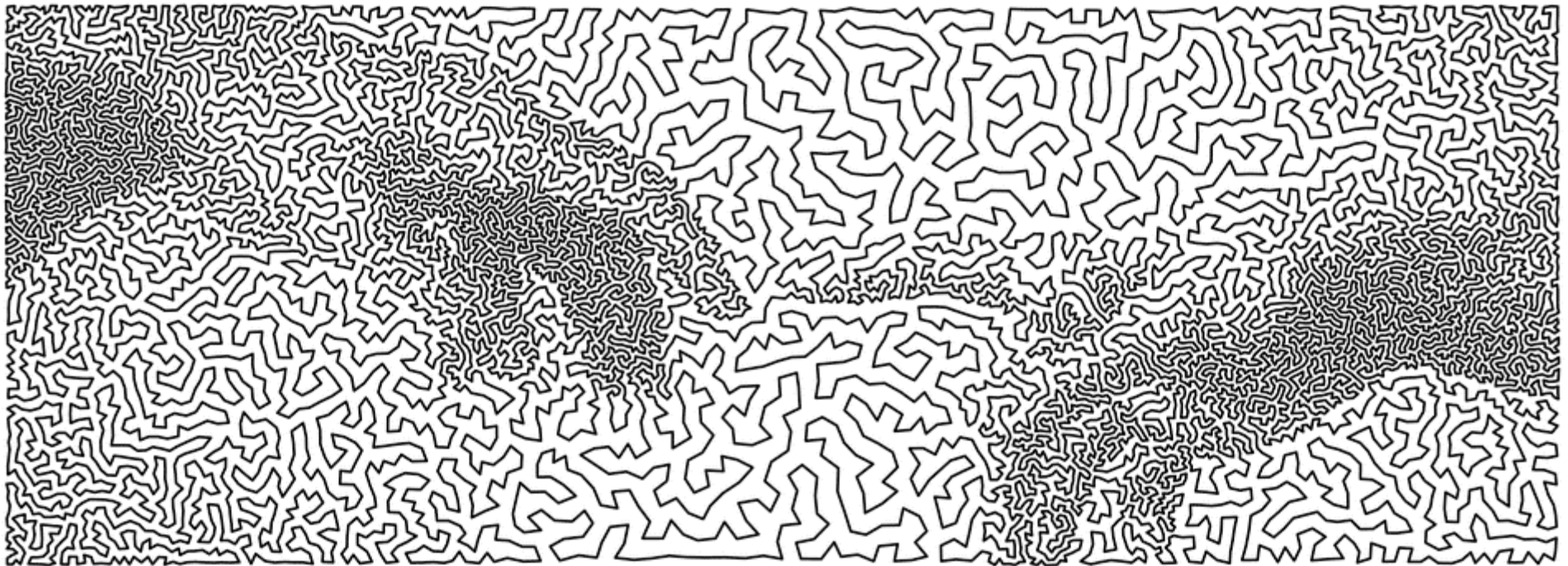


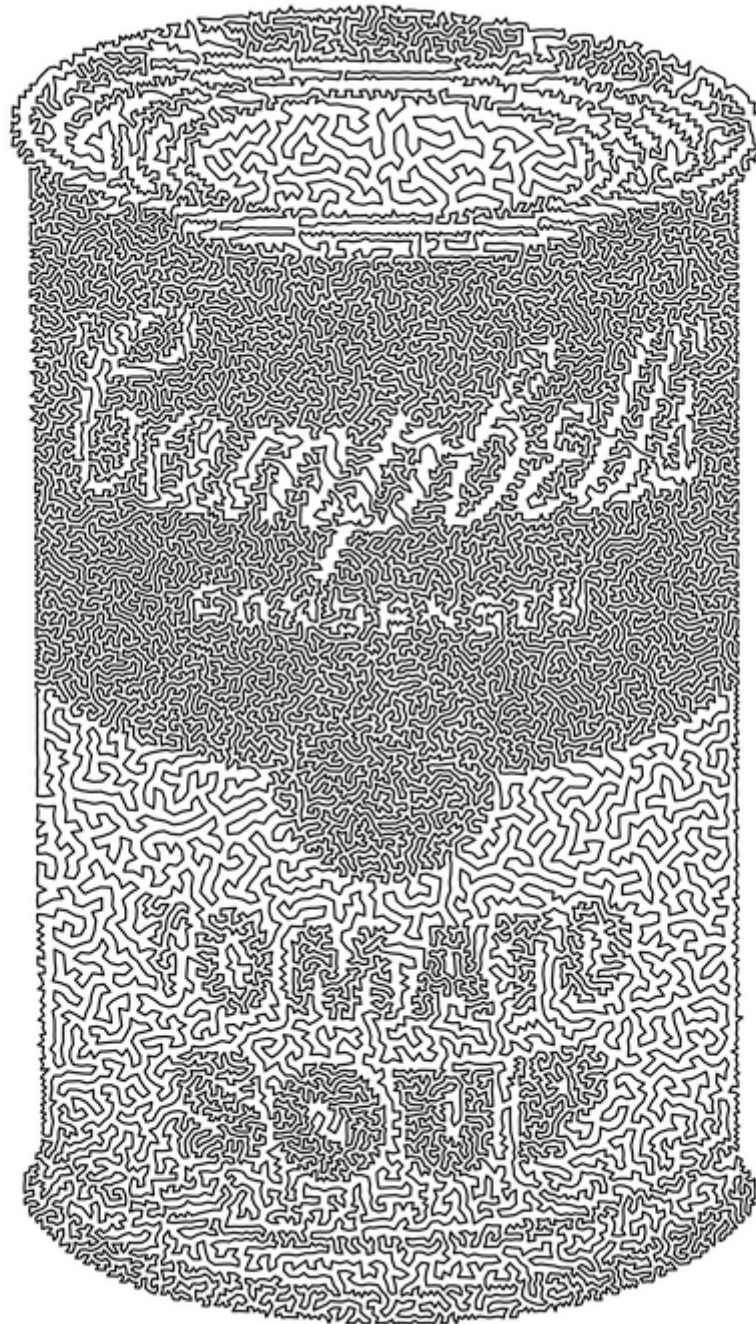


Hands

INFORMS 06

Pittsburgh, PA



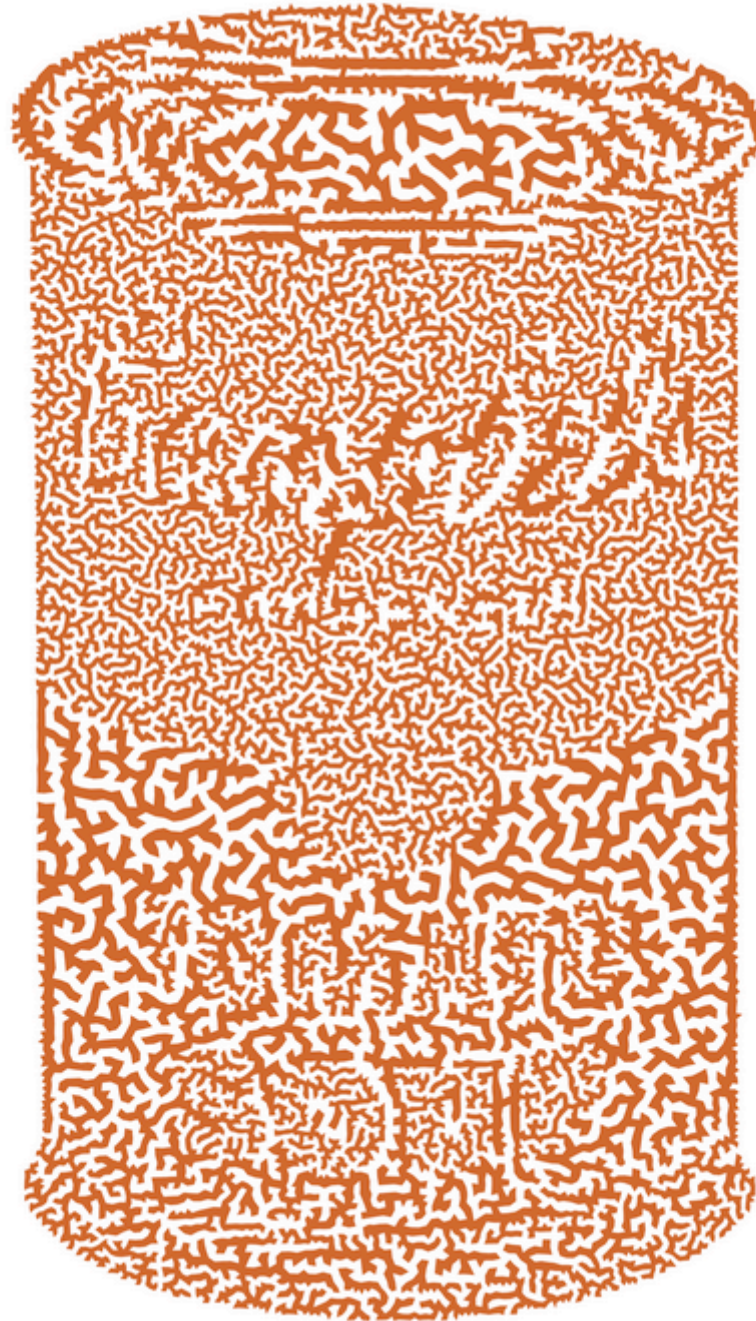


What's Inside?

INFORMS 06
Pittsburgh PA

JMM 2007
San Diego CA

Bridges 2007
Donostia, Spain



This is!

INFORMS 06
Pittsburgh PA

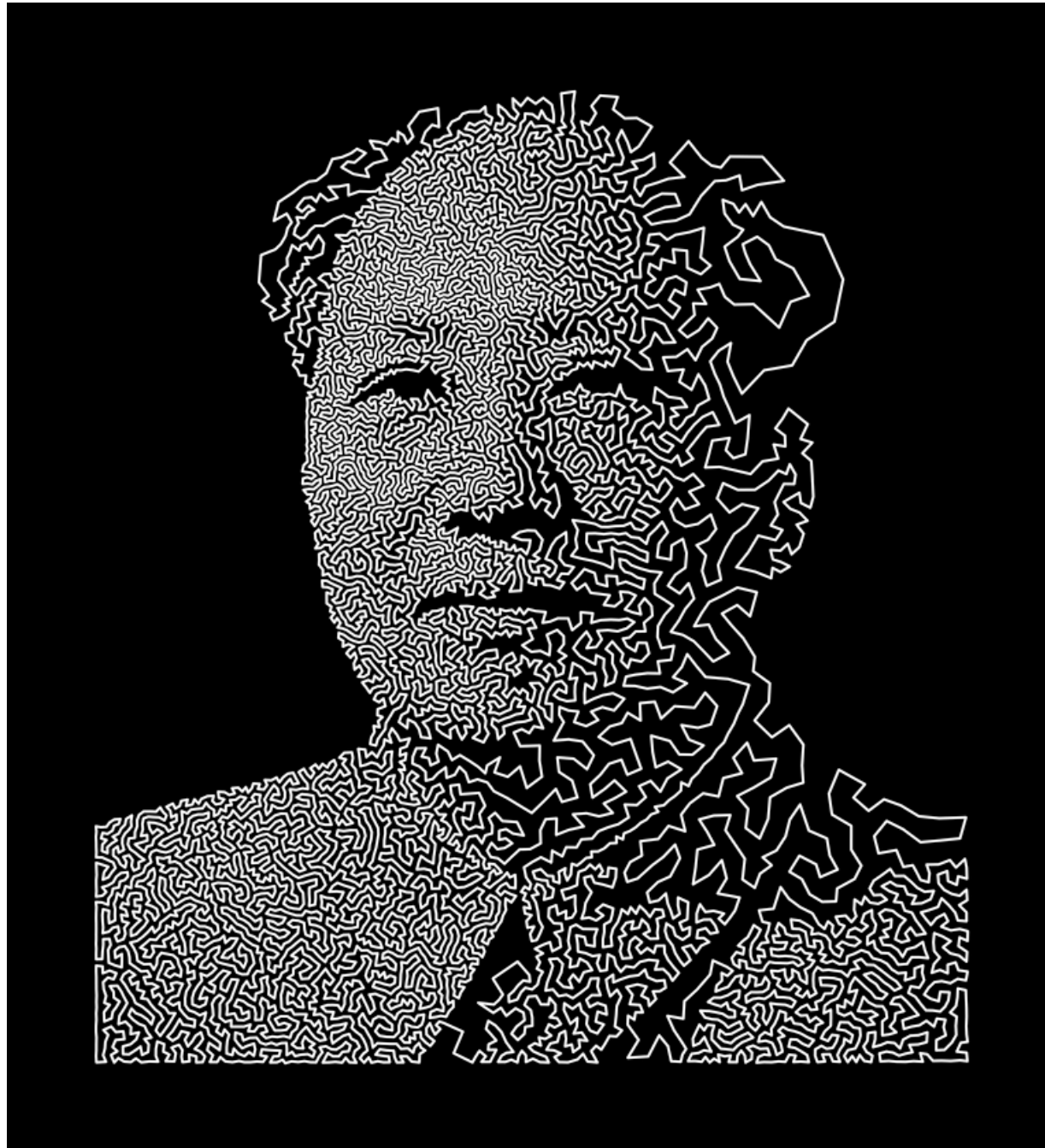
JMM 2007
San Diego CA

Bridges 2007
Donostia, Spain



One loop variation 1

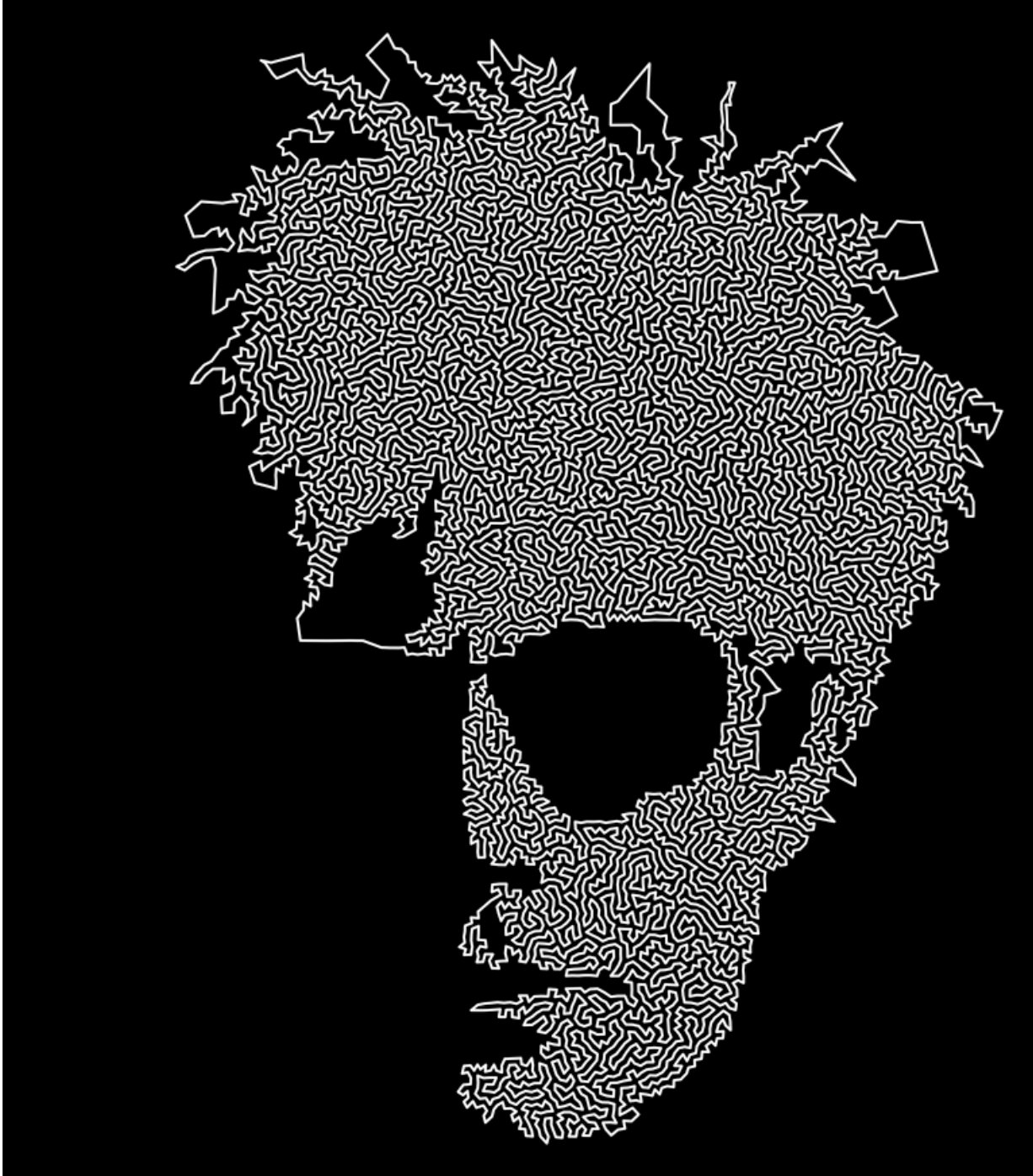
INFORMS 06
Pittsburgh PA



One loop variation 2

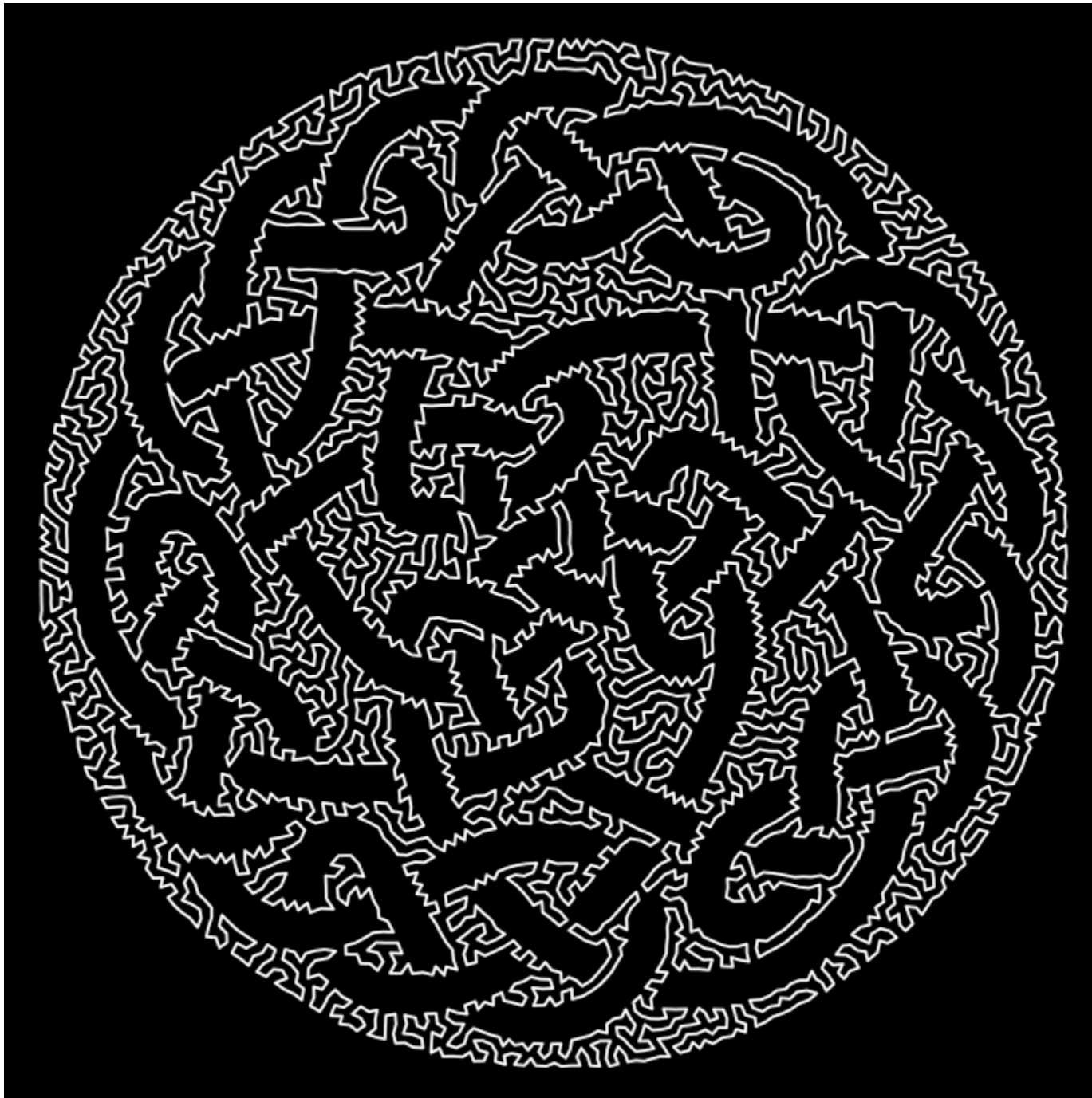
INFORMS 06

Pittsburgh PA



One loop variation 3

INFORMS 06
Pittsburgh PA



Knot?

CPAIOR06
Cork, Ireland

JMM 2007
New Orleans, LA

Bridges 2007
Donostia, Spain

One fish, two fish, red fish, black fish

JMM 2008
San Diego, CA

Bridges 2008
Leeuwarden, NL

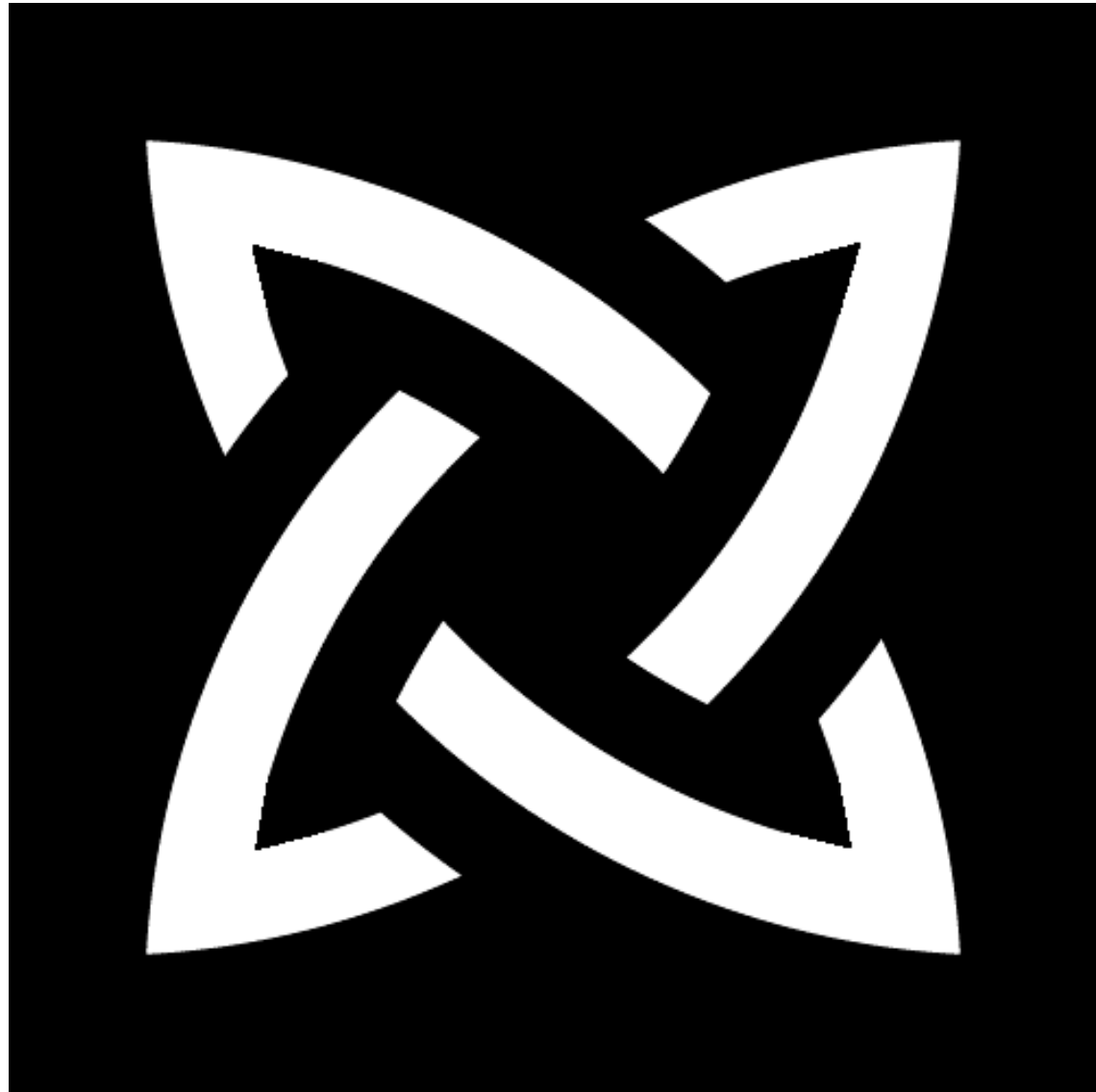


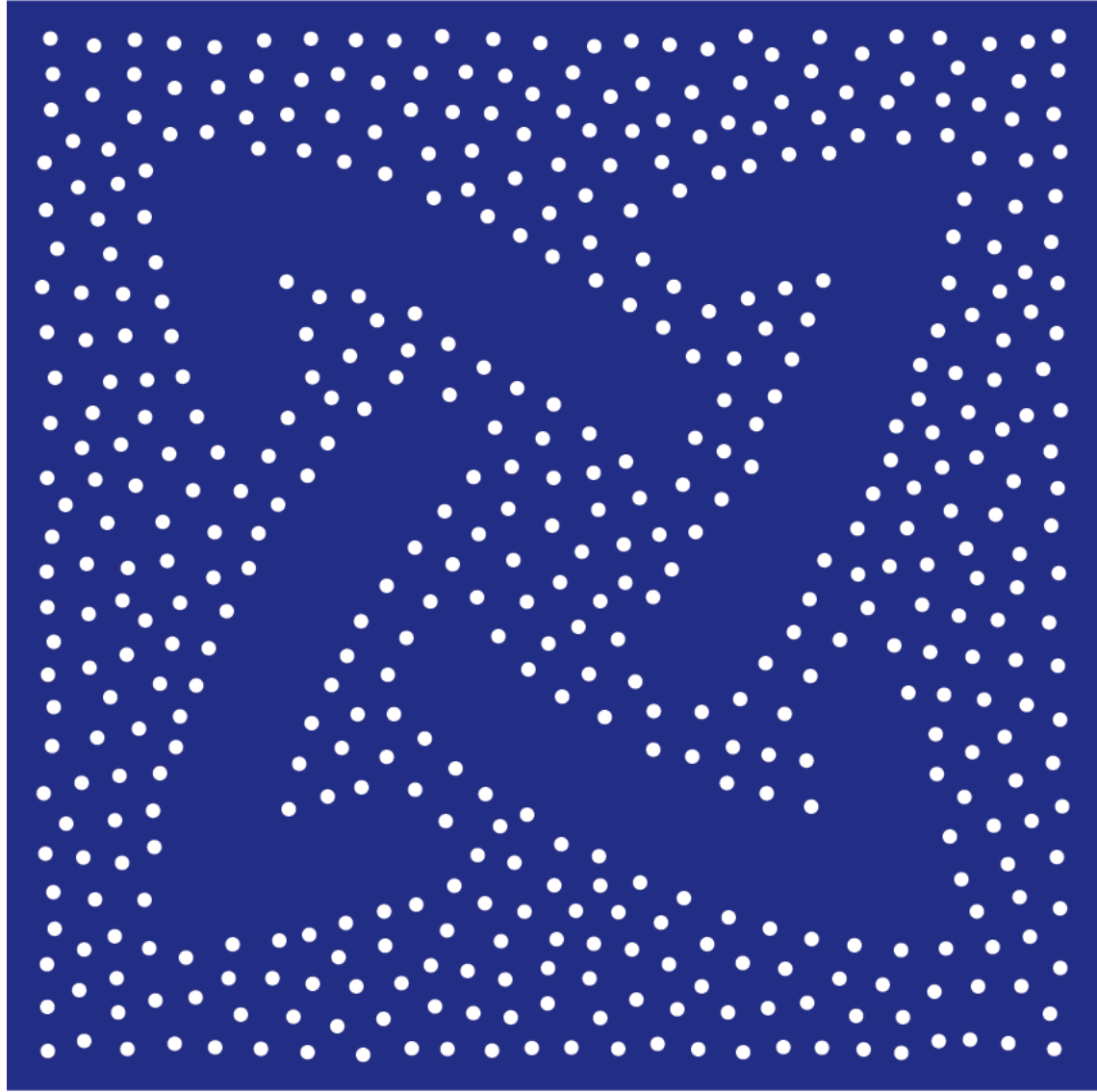


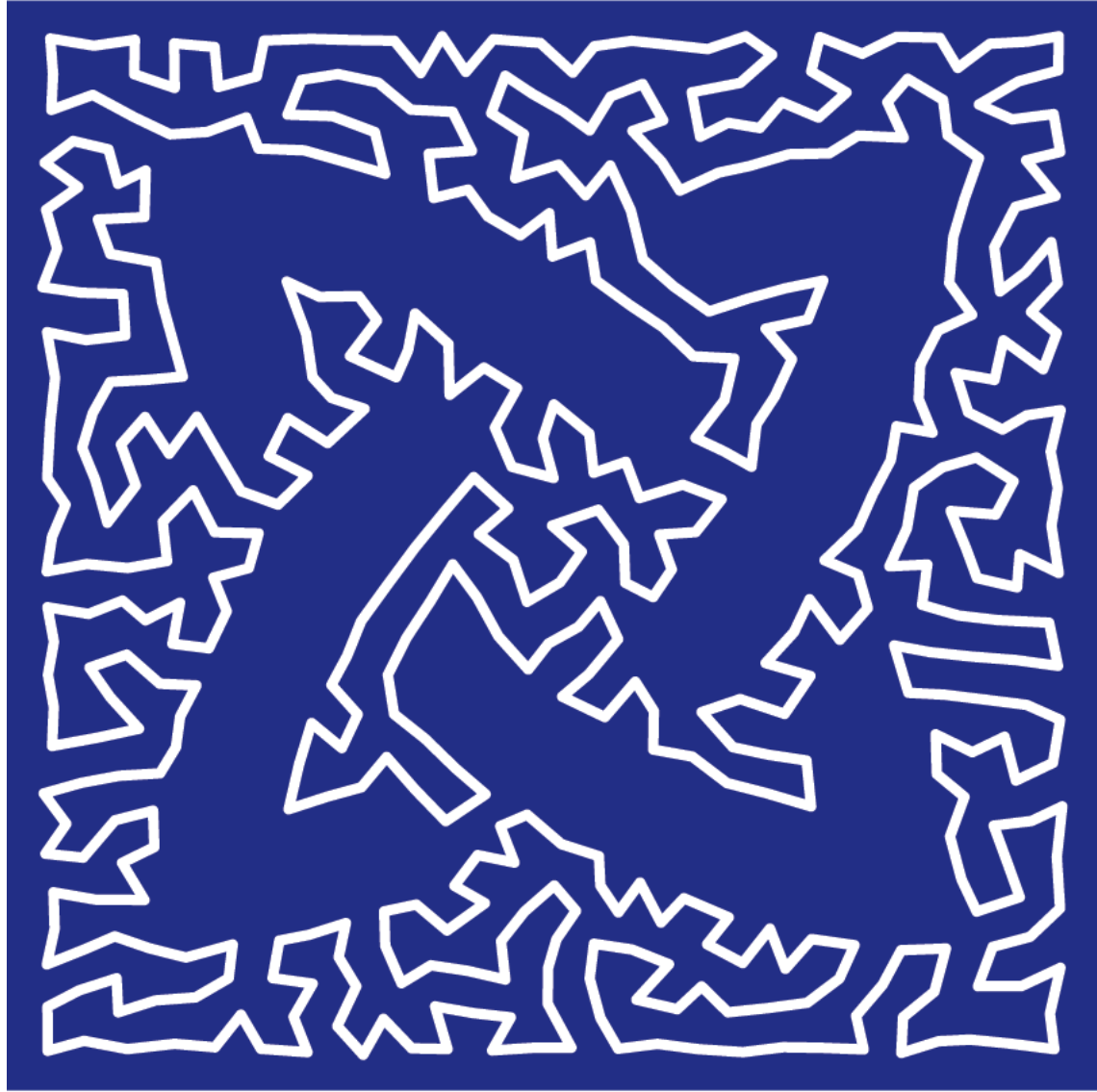
Outside Ring

JMM 2008
San Diego, CA

Bridges 2008
Leeuwarden, NL

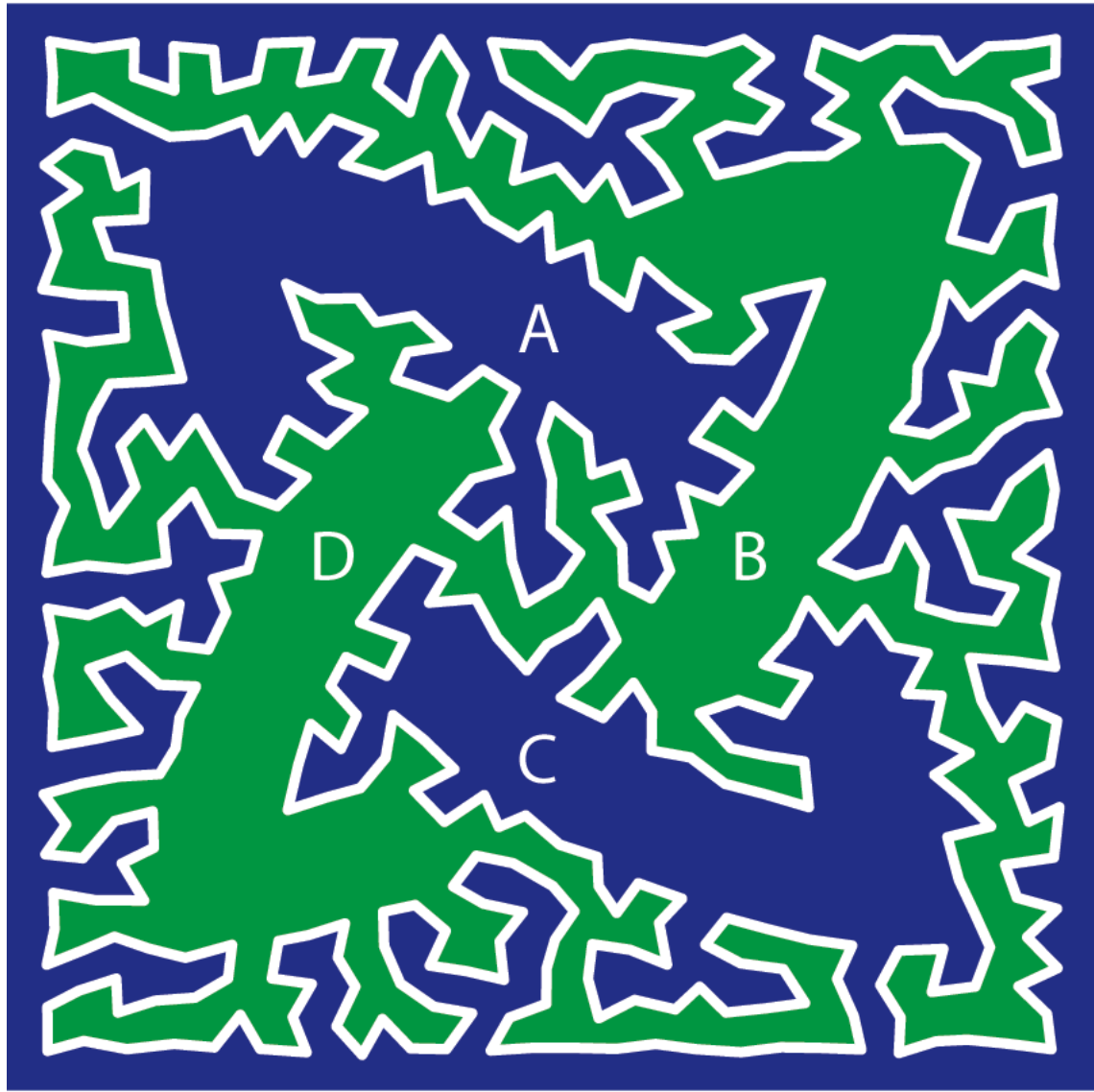


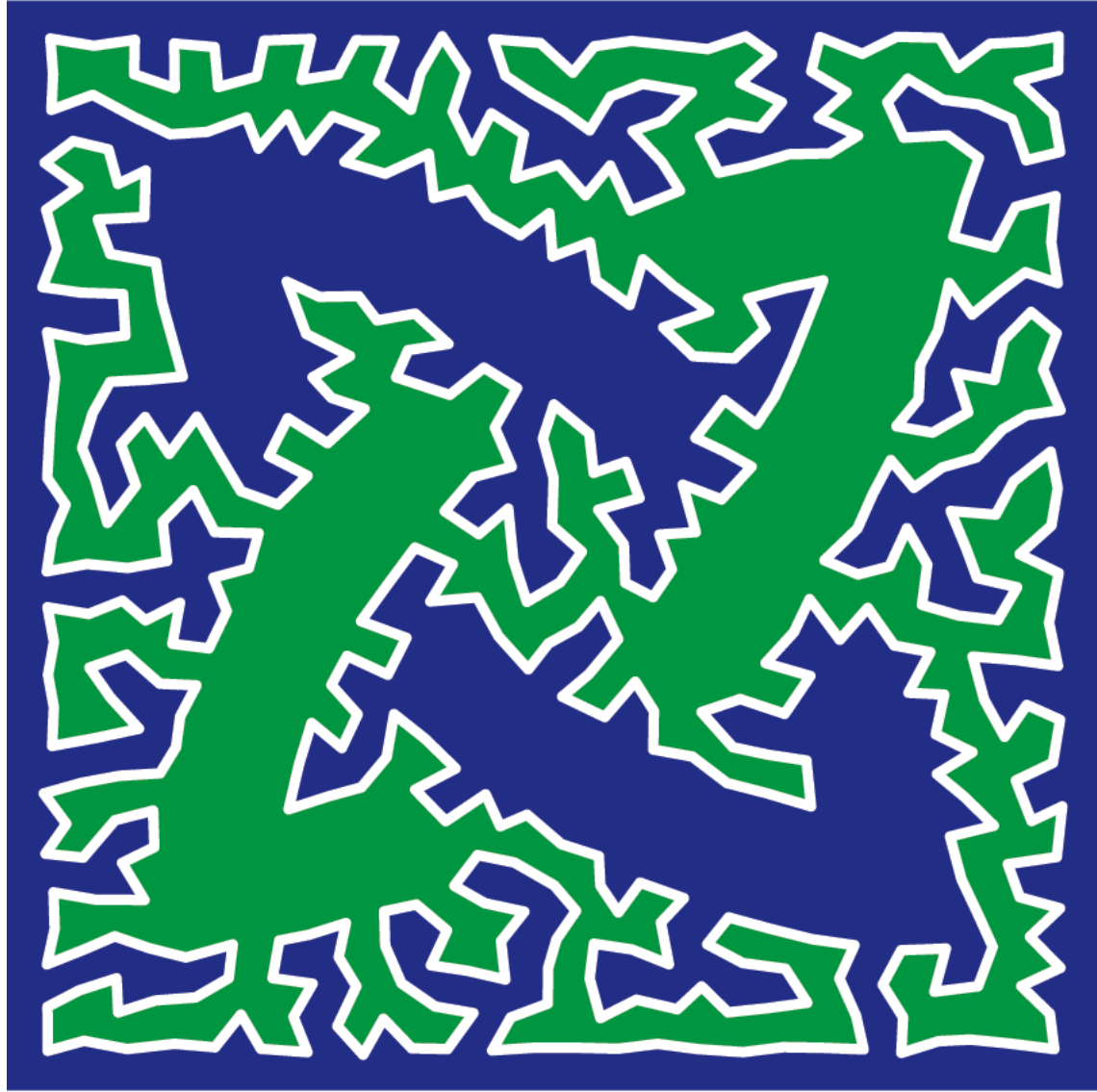








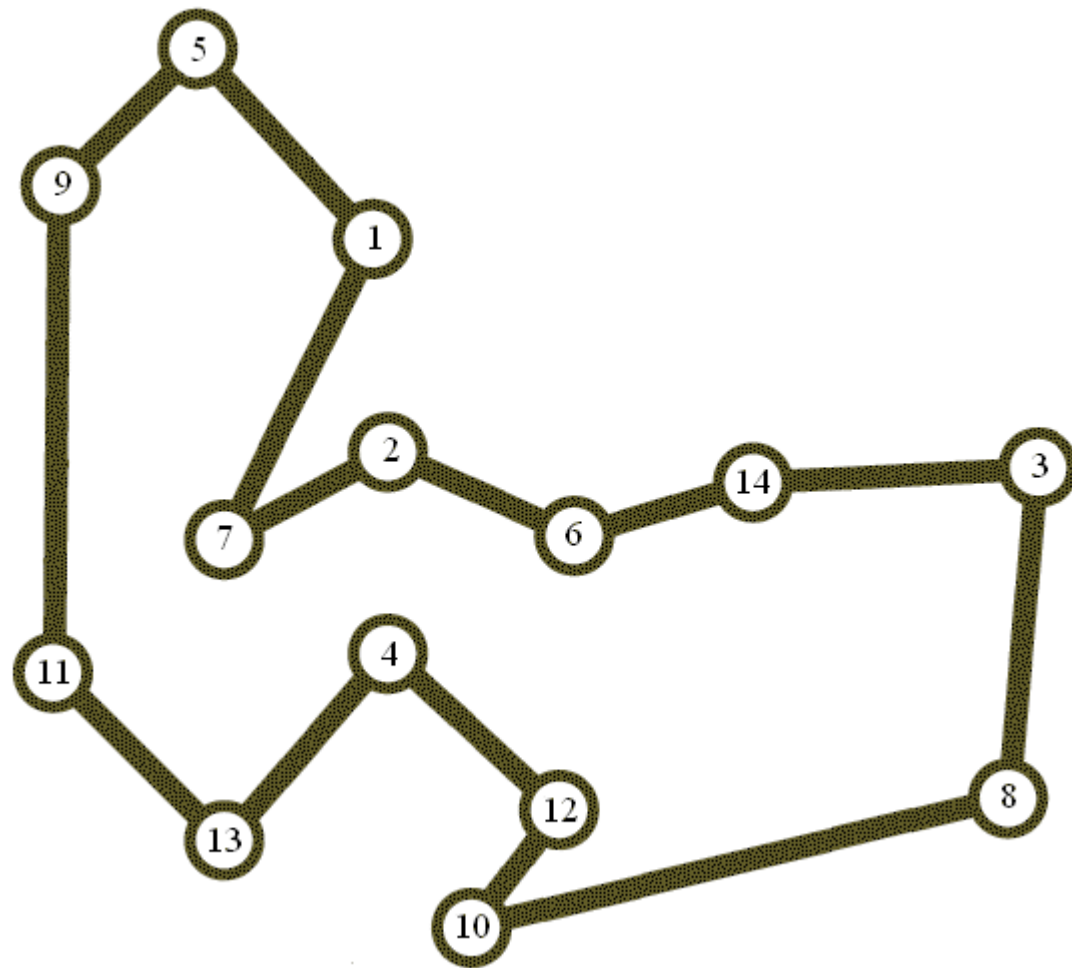




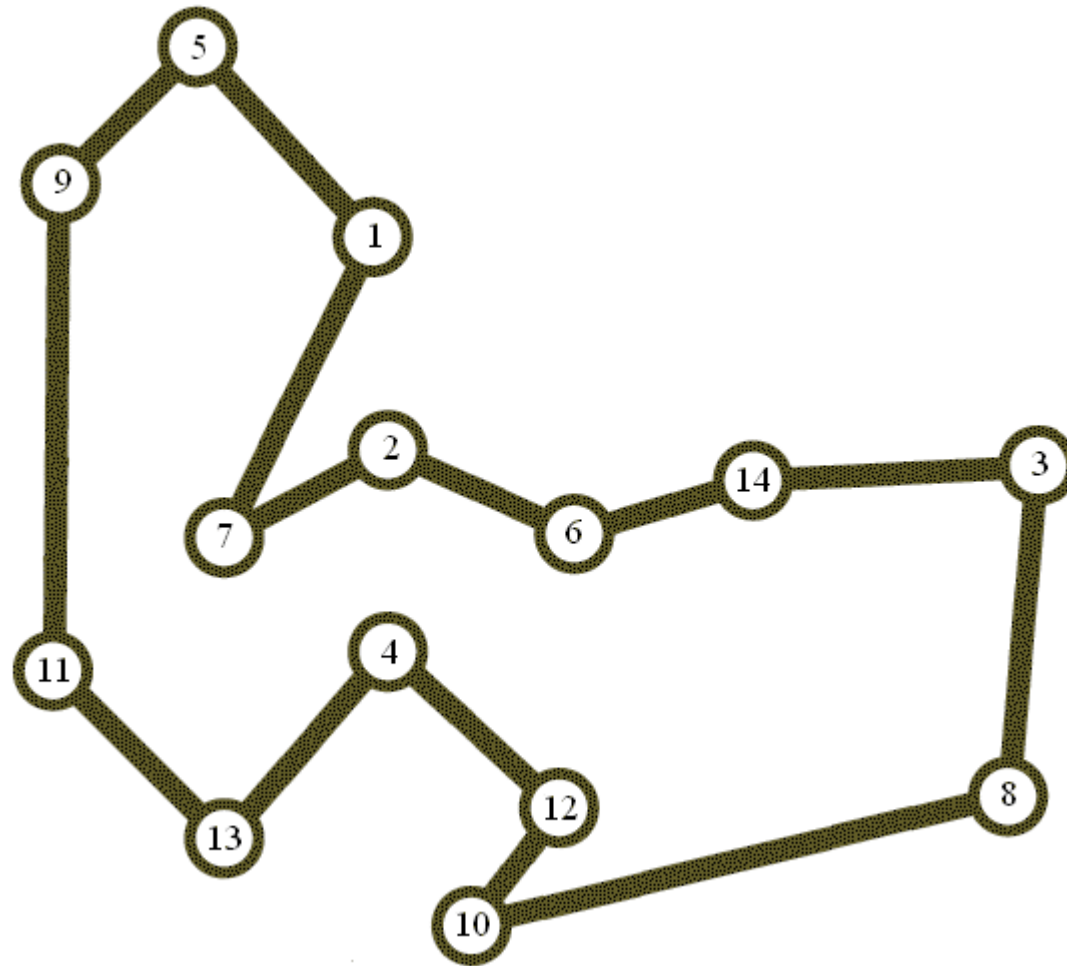
Solving TSPs with IP:



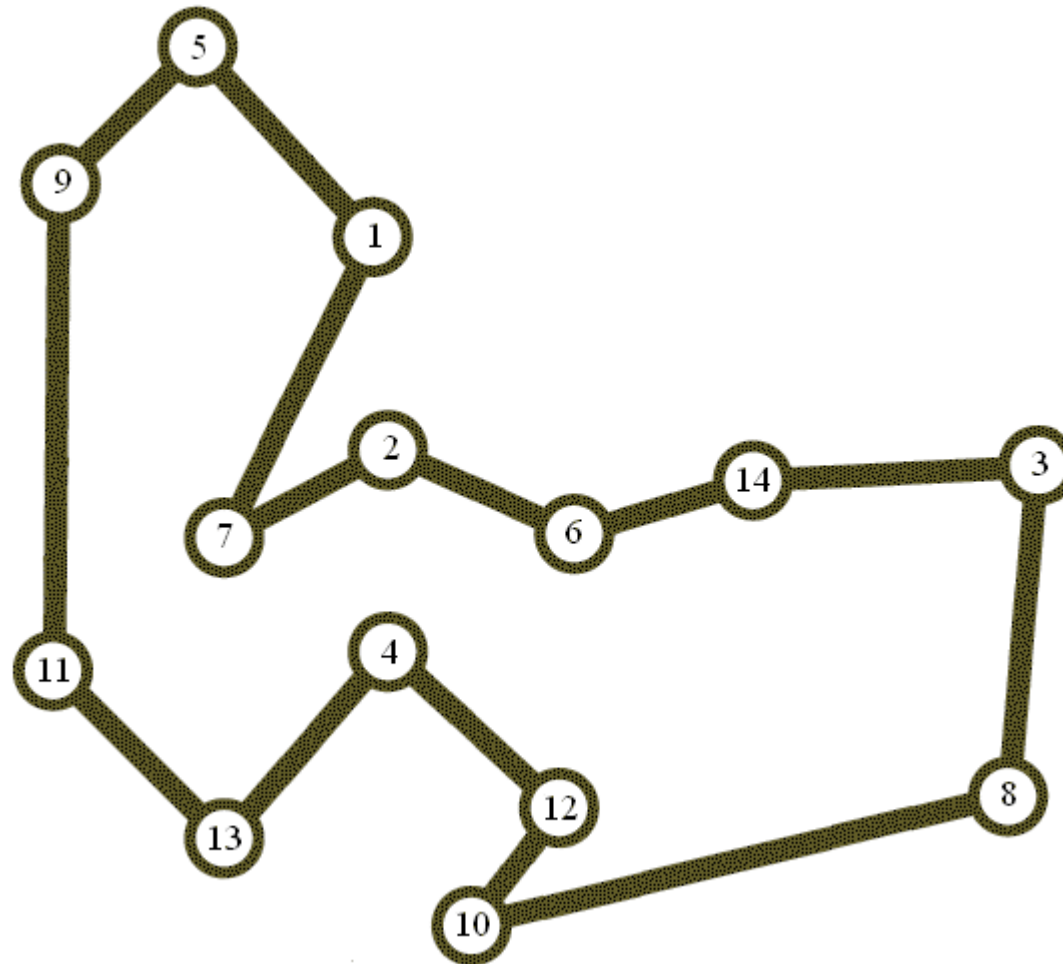
Eighth Judicial Circuit traveled by Lincoln in 1850.



Variables: $x_e = 1$ if edge e is used in the tour



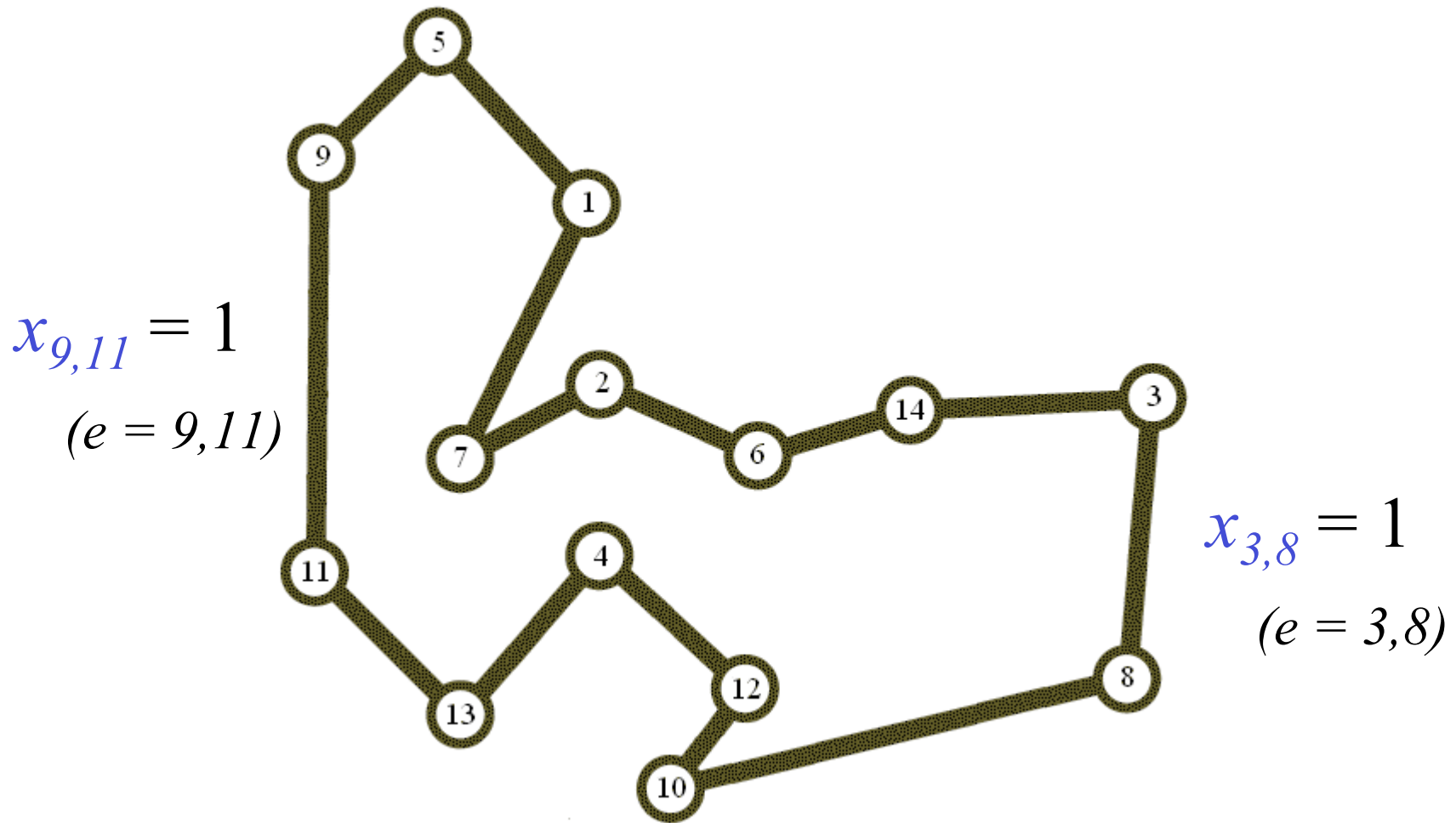
Variables: $x_e = 1$ if edge e is used in the tour



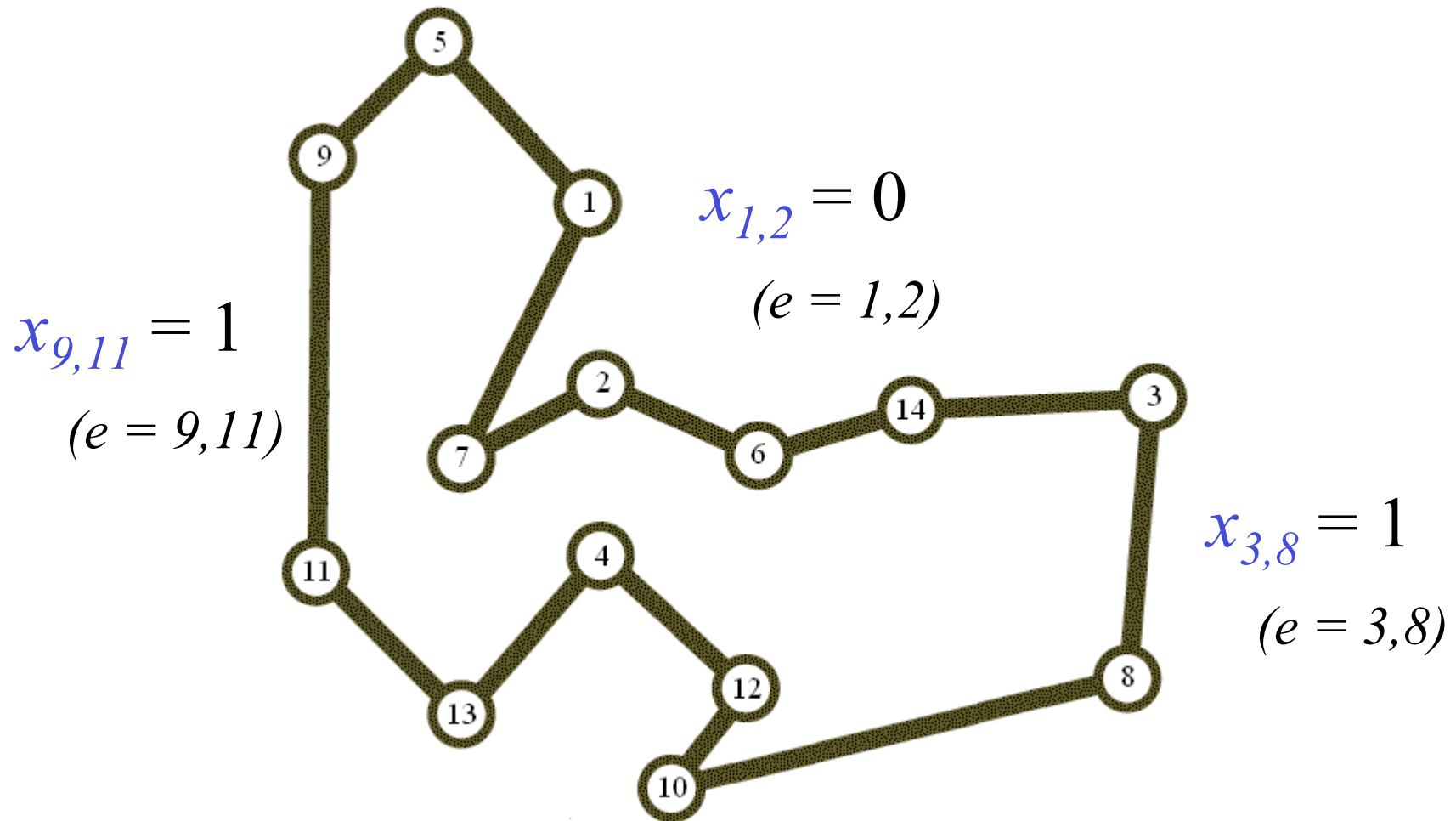
$$x_{3,8} = 1$$

($e = 3,8$)

Variables: $x_e = 1$ if edge e is used in the tour

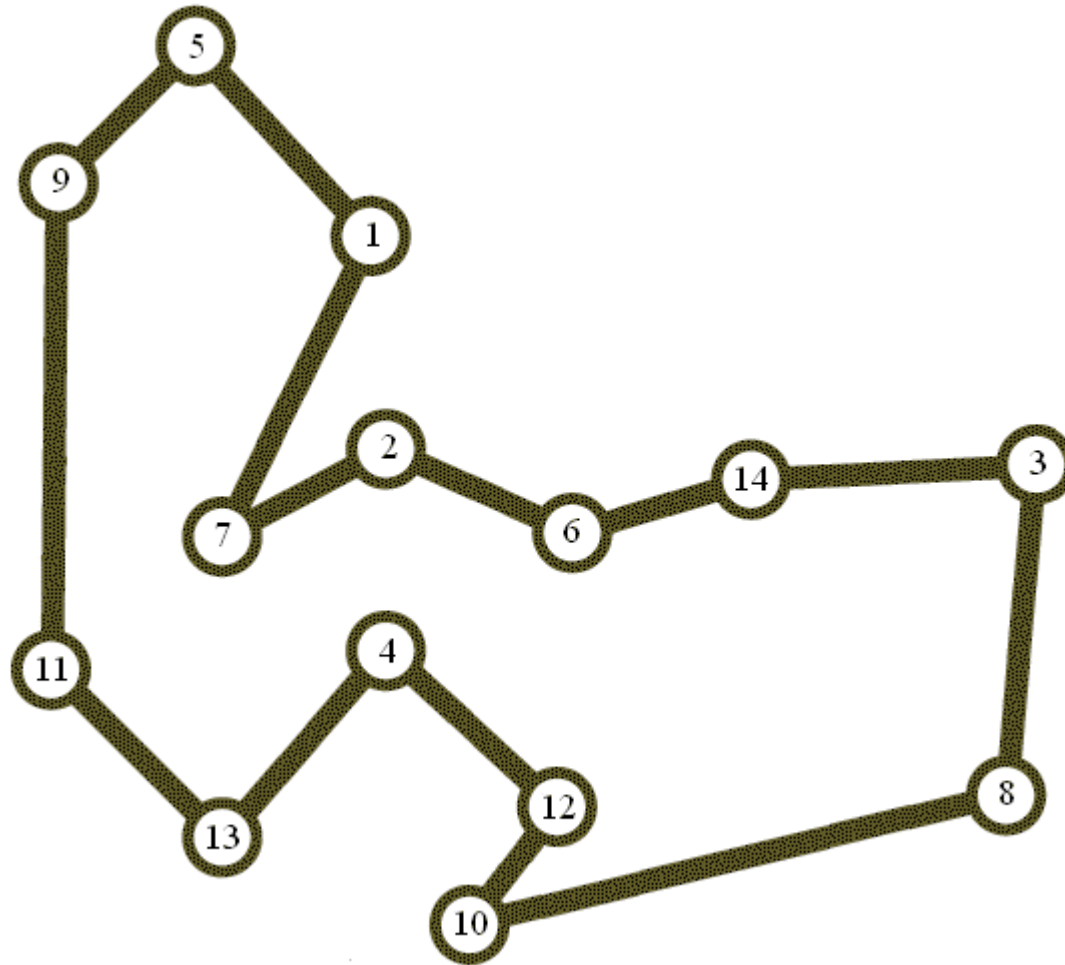


Variables: $x_e = 1$ if edge e is used in the tour



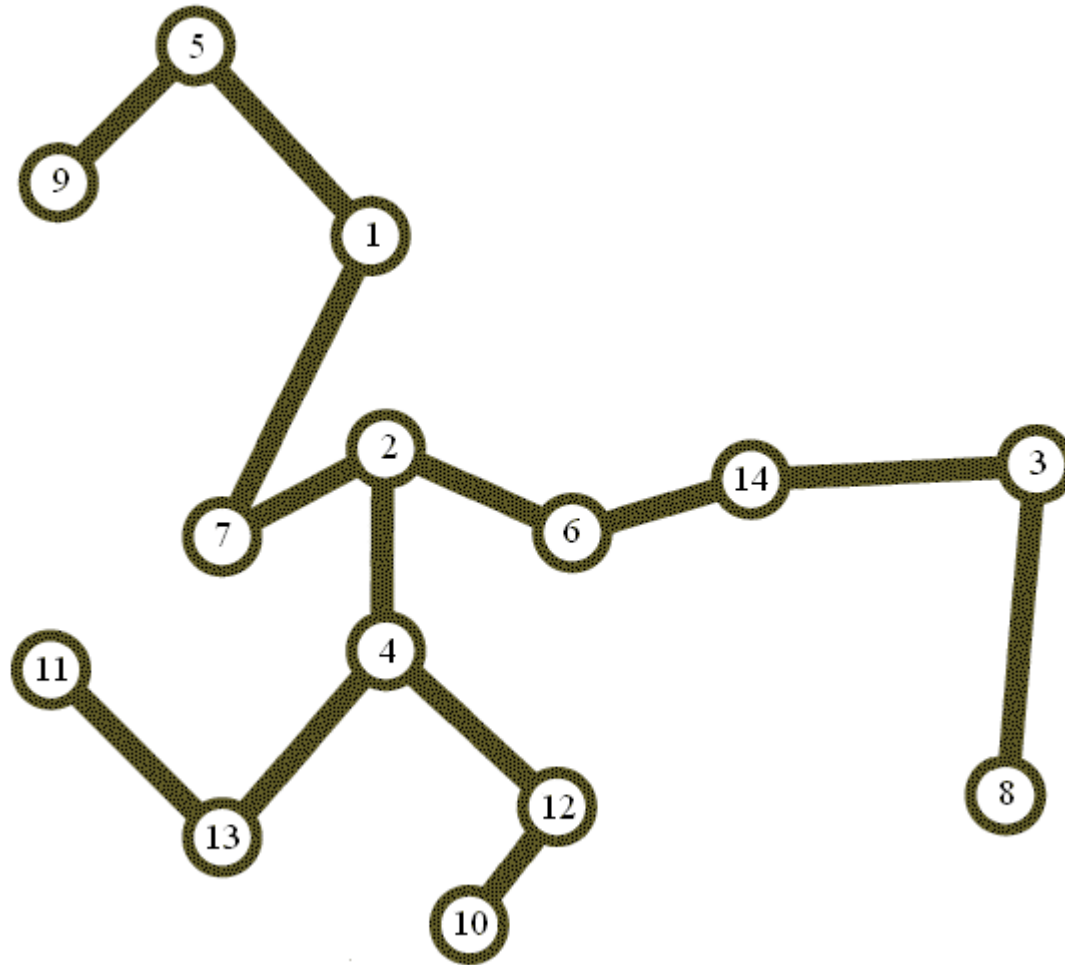
Objective: minimize length of tour

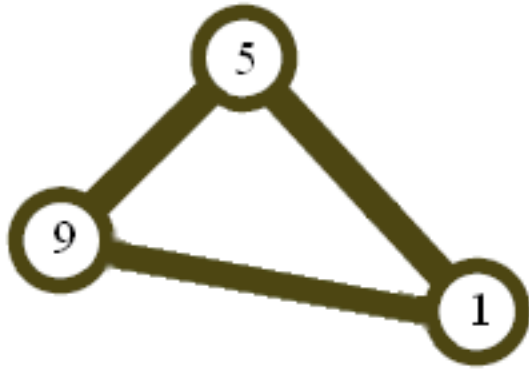
$$\text{total cost} = \sum (c_e x_e : \text{all edges } e)$$



Constraints: each city v must be touched twice

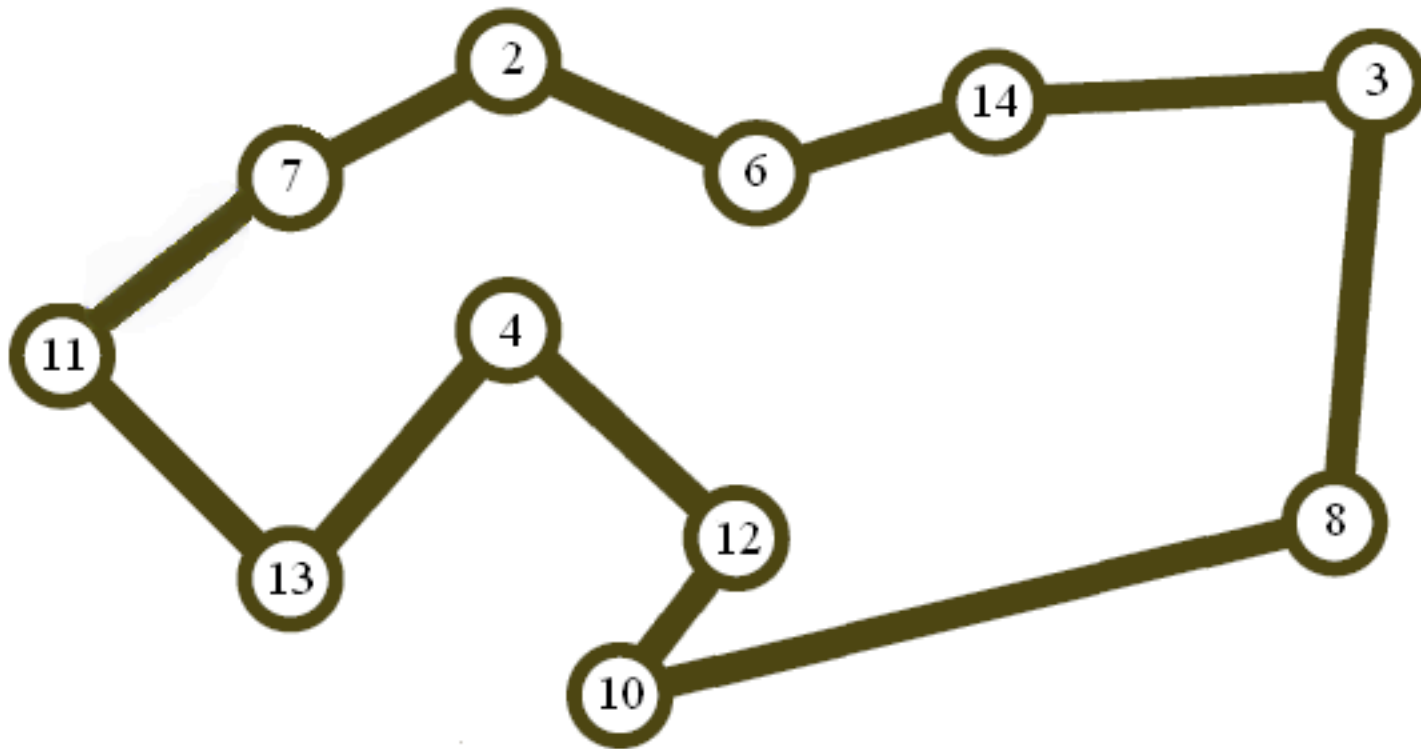
$$\sum (x_e : \text{all edges } e \text{ that touch } v) = 2$$





Constraints: no subtours!

$$\sum (x_e : \text{all } e \text{ with both endpoints in } S) < |S|$$



Solving TSPs with IP (DFJ formulation):

Variables: $x_e = 1$ if edge e is used in the tour

Objective: minimize length of tour

$$\text{total cost} = \sum (c_e x_e : \text{all edges } e)$$

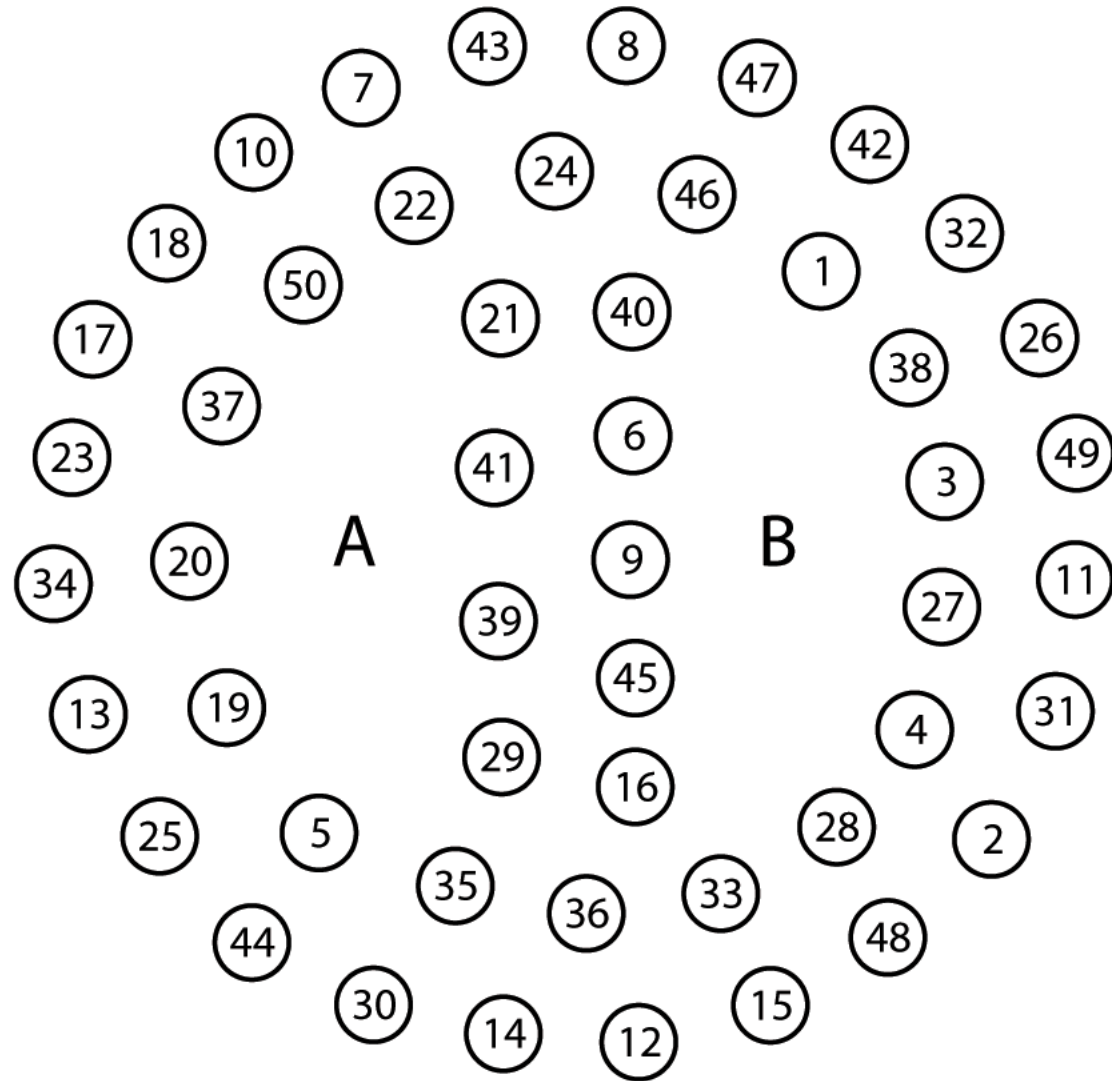
Constraints: each city v must be touched twice

$$\sum (x_e : \text{all edges } e \text{ that touch } v) = 2 \quad \forall v$$

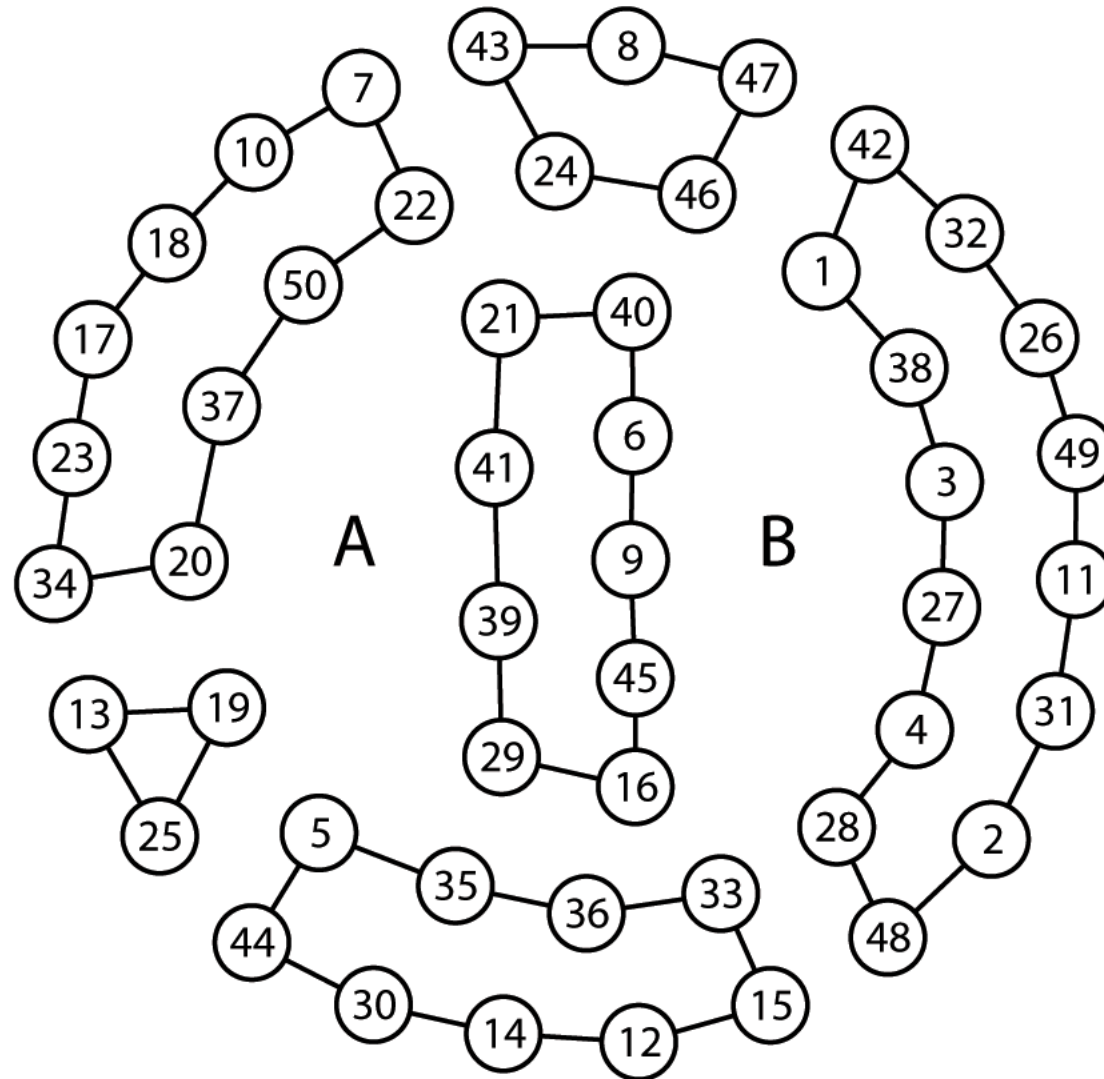
Constraints: no subtours are allowed

$$\sum (x_e : \text{all } e \text{ w/both endpts in } S) < |S| \quad \forall S$$

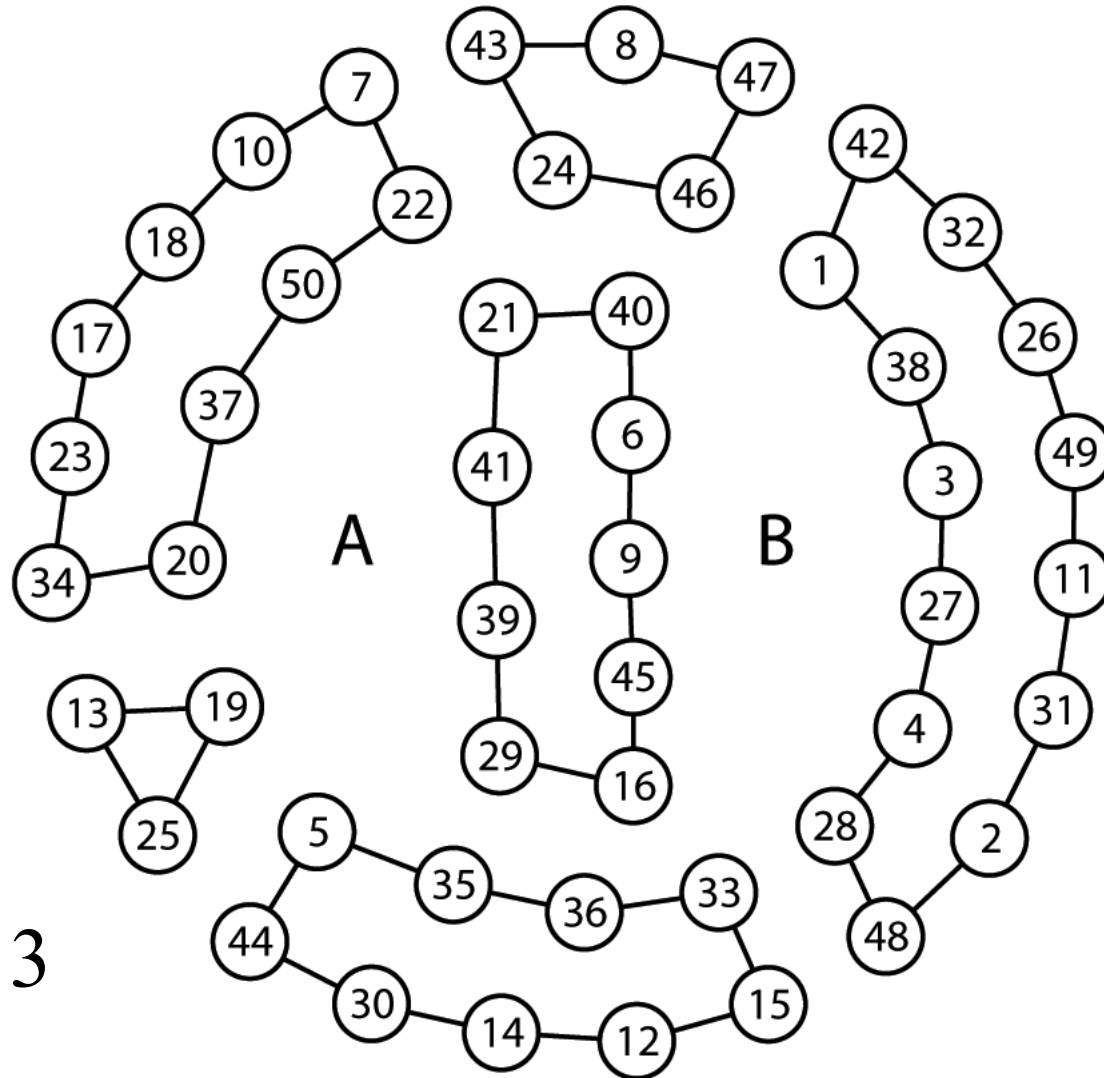
An example:



Stage 1: 6 subtours.

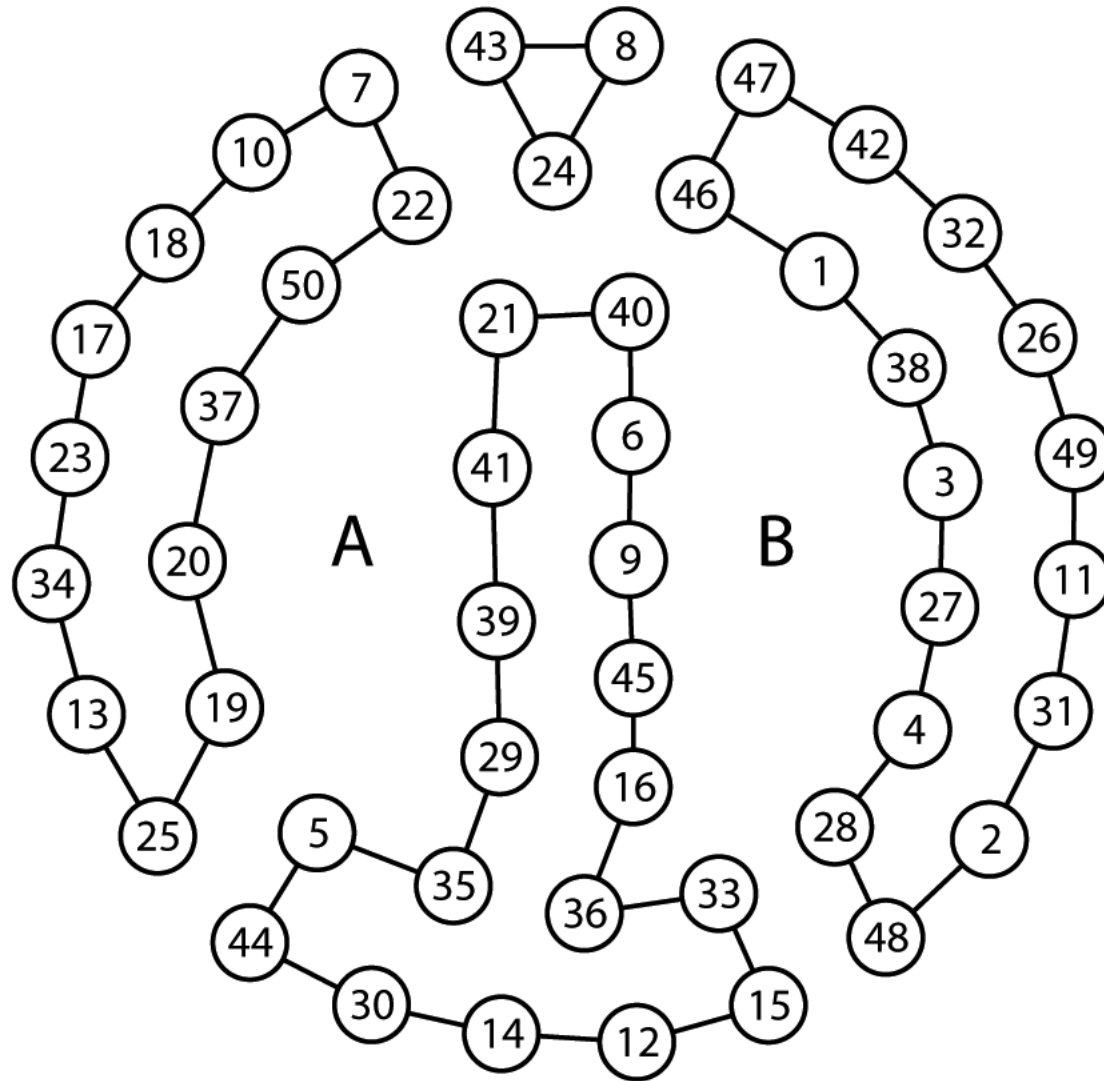


Stage 1: 6 subtours.

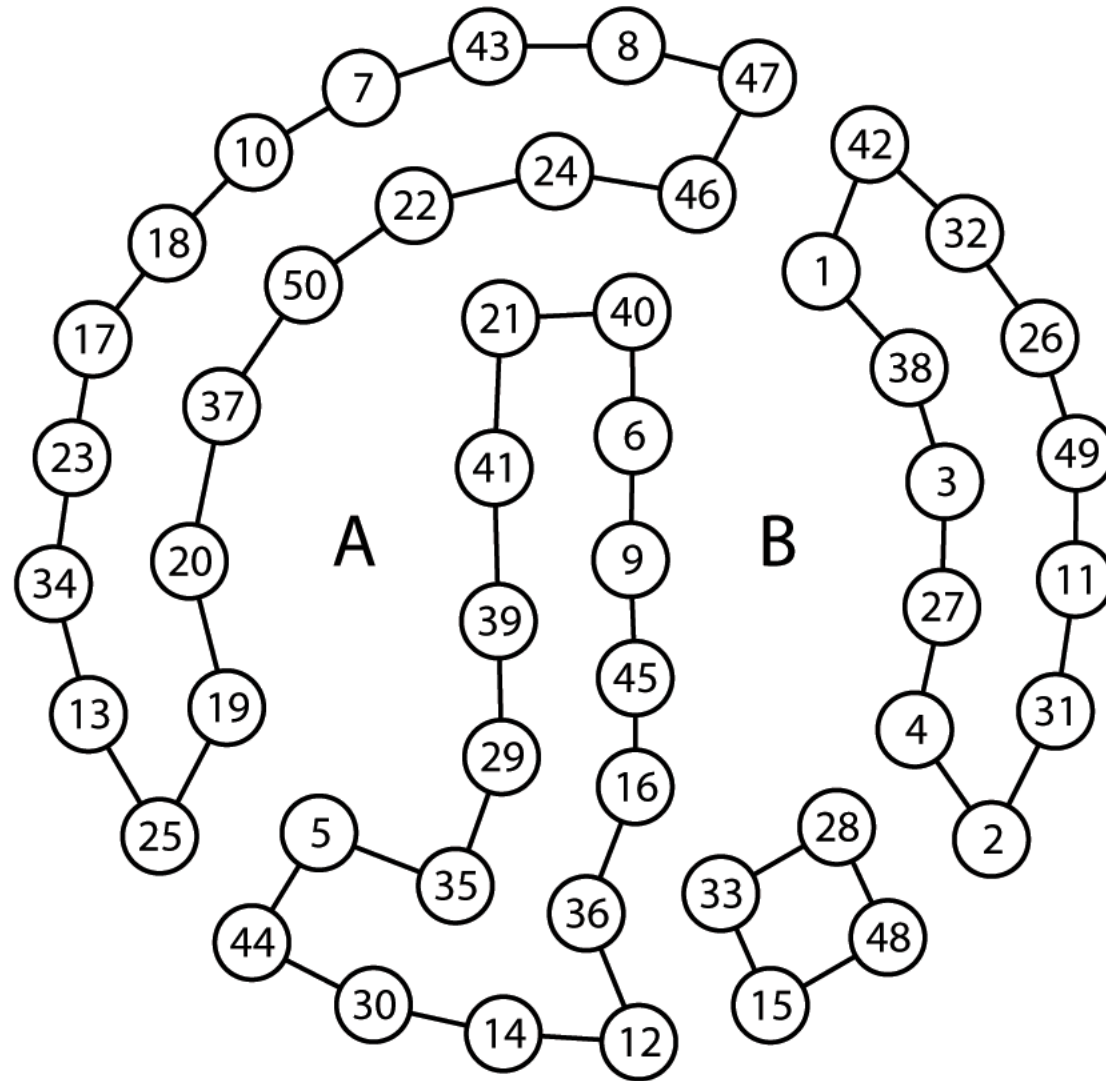


$$\begin{aligned} & x_{13,19} \\ + & x_{13,25} \\ + & x_{19,25} < 3 \end{aligned}$$

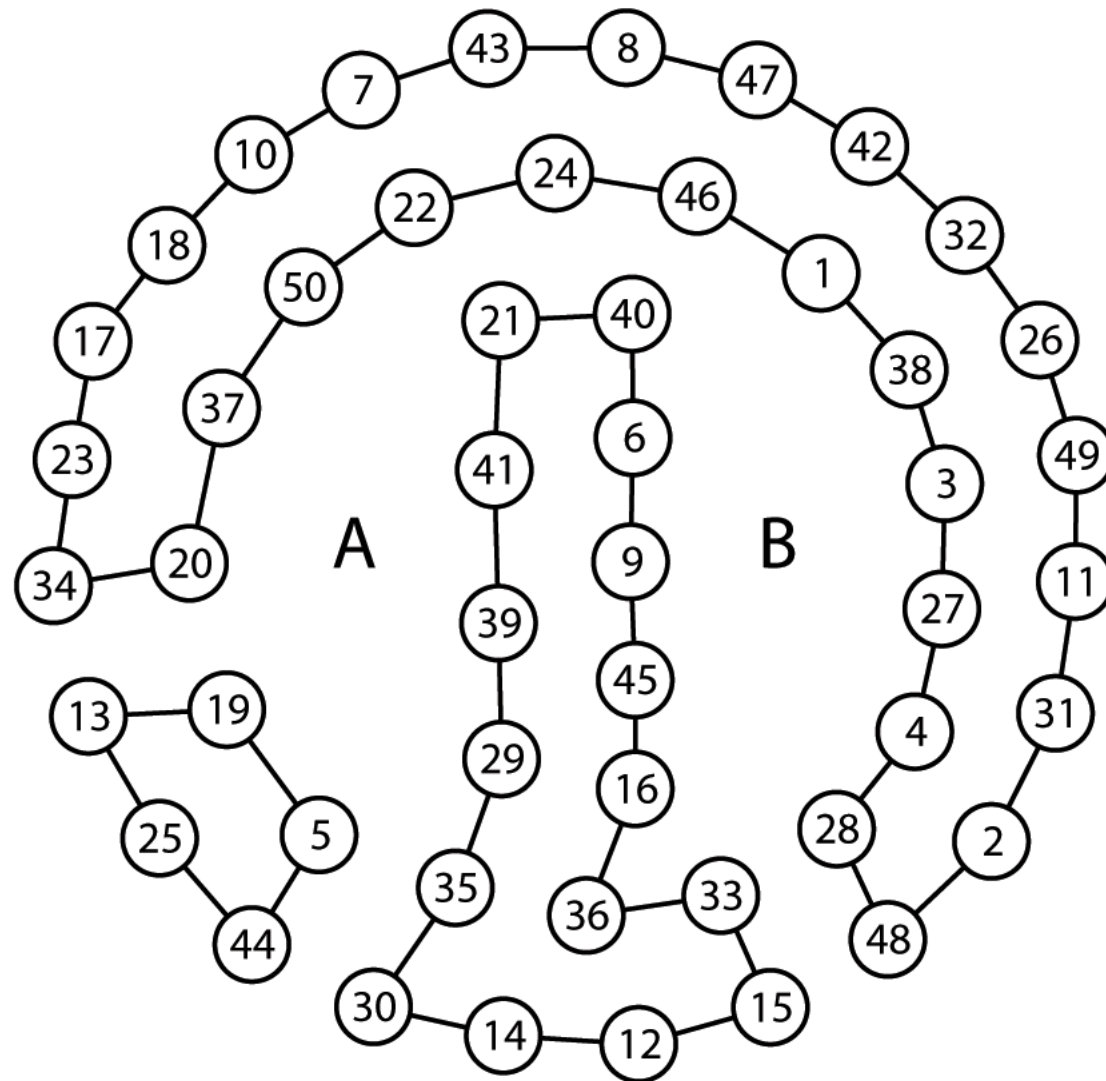
Stage 2: 4 more subtours.



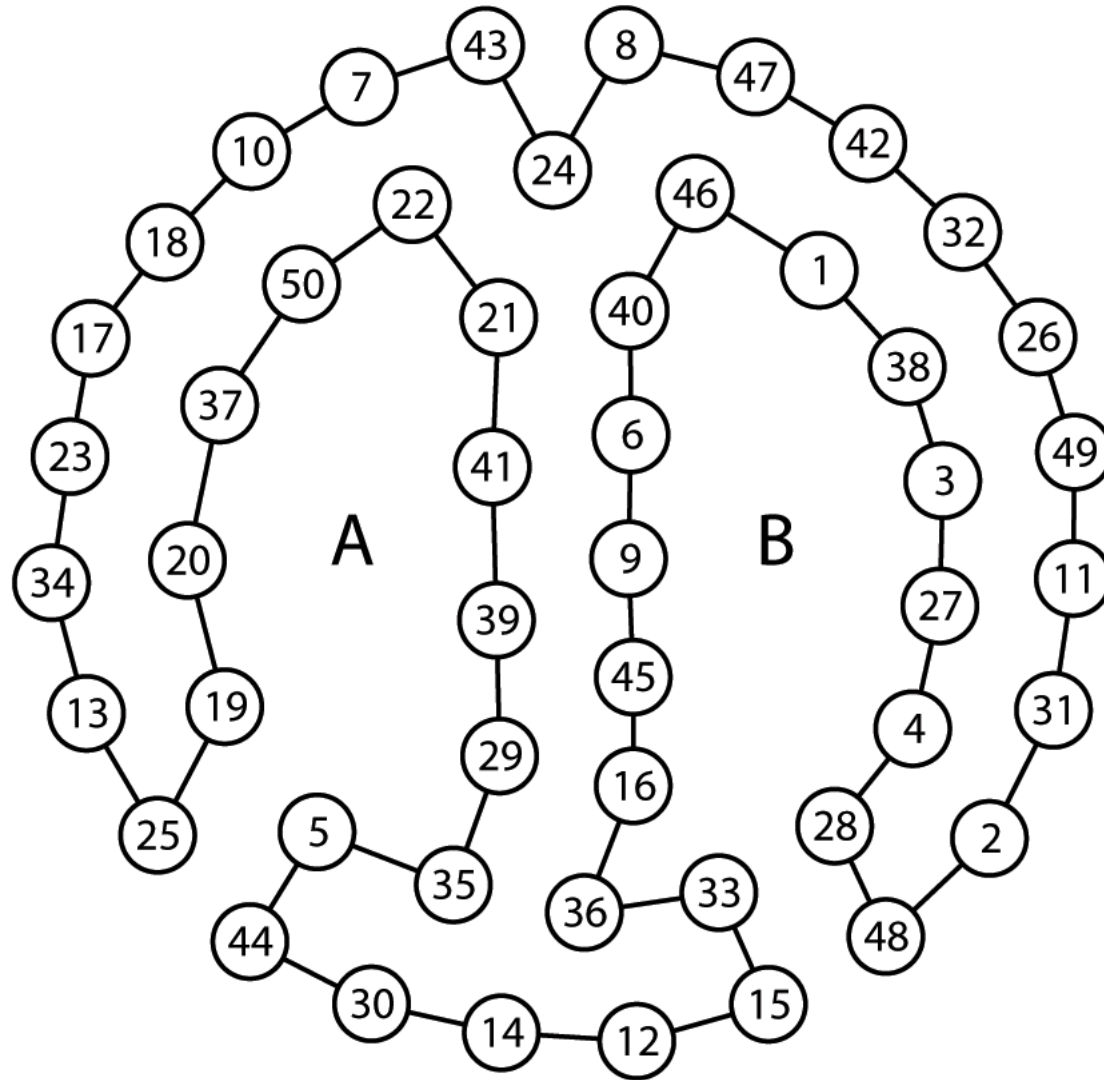
Stage 3: 4 more subtours.

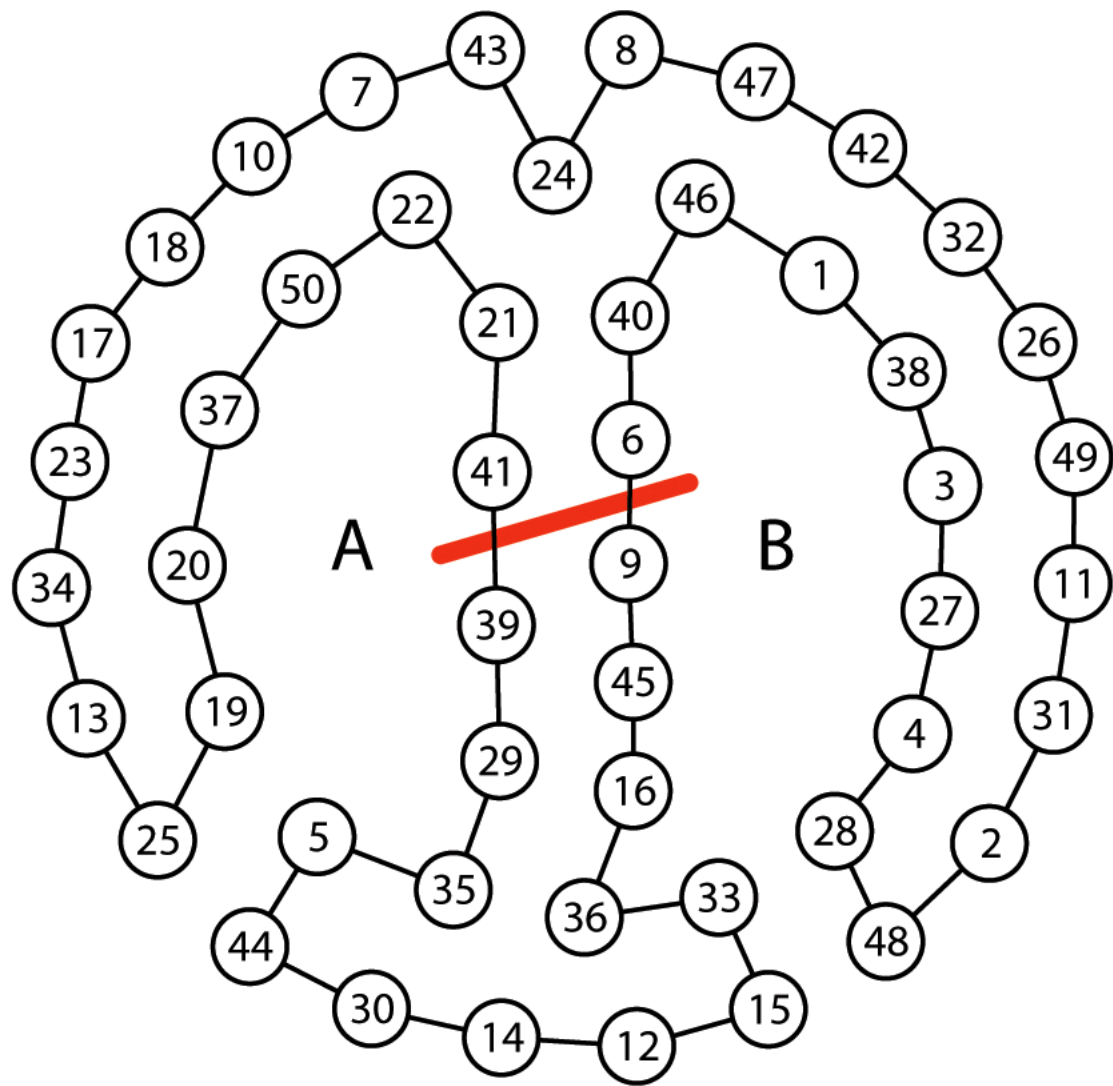


Stage 4: 3 more subtours.



Stage 5: success!





Definition:

$$\chi(l_{A,B}) = \{ \text{all edges } e \text{ that cross line } l_{A,B} \}$$

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$$\chi(l_{A,B}) = \{ \text{all edges } e \text{ that cross line } l_{A,B} \}$$

Note:

A and B lie on the same side of the tour



$\Sigma (x_e : \text{all edges } e \text{ in } \chi(l_{A,B}))$ is even



$$\Sigma (x_e : \text{all edges } e \text{ in } \chi(l_{A,B})) = 2 y_{A,B}$$

Definition:

$$\chi(l_{A,B}) = \{ \text{all edges } e \text{ that cross line } l_{A,B} \}$$

Equivalently:

A and B lie on opposite sides of the tour



$\Sigma (x_e : \text{all edges } e \text{ in } \chi(l_{A,B}))$ is odd



$$\Sigma (x_e : \text{all edges } e \text{ in } \chi(l_{A,B})) = 2 y_{A,B} + 1$$

The 'Big' Idea:

Use

$$\sum (x_e : \text{all edges } e \text{ in } \chi(l_{A,B})) = 2 y_{A,B}$$

and/or

$$\sum (x_e : \text{all edges } e \text{ in } \chi(l_{A,B})) = 2 y_{A,B} + 1$$

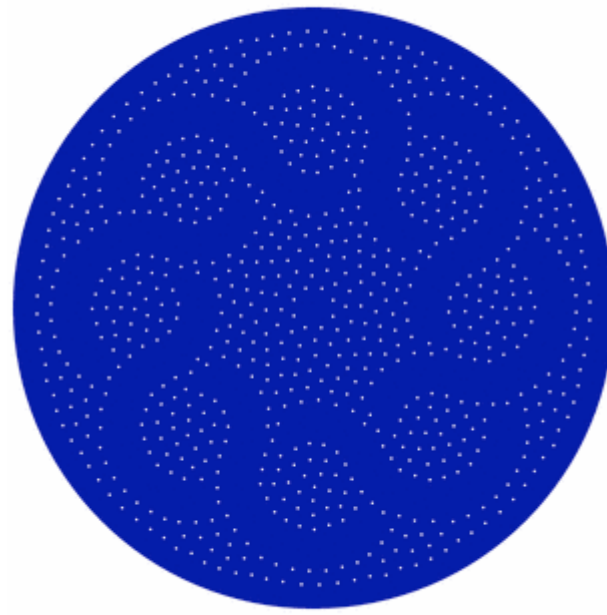
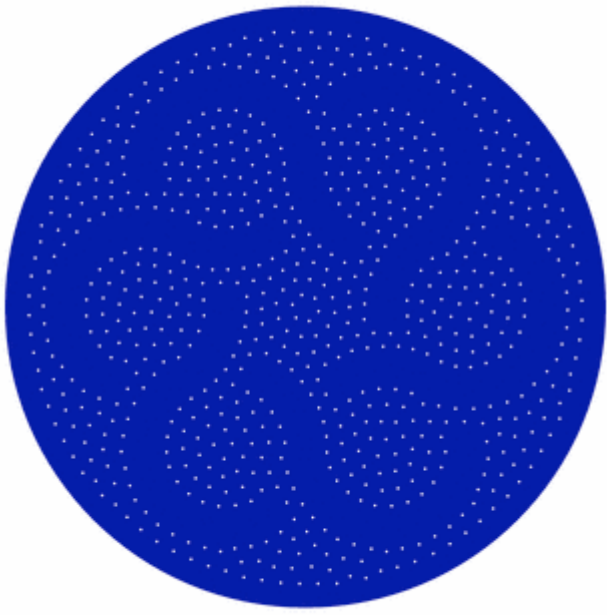
in the DFJ IP formulation.





































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- C.S. Kaplan and R. Bosch. 2005. TSP art. *Bridges 2005*.
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